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과제명: IoT 환경을 위한 고성능 플래시 메모리 스토리지 기반 인메모리 분산 DBMS 연구개발

SW STOR LAB

Software Technology Advanced Research

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## Introduction

- Changes in Trends
  - Device .
    - ✓ Demand for rapid processing of large volumes of data
      - → Need high-performance device



Hard Disk Drive & Solid-State Drive



**Dynamic Random-Access Memory** 

- ✓ DRAM Cost per GB(\$/GB) is expensive
  - ✓ DRAM Cost per GB: \$2.57
  - ✓ SSD Cost per GB: \$0.23









## Introduction



- In-Memory Key-Value Stores require memory space for system operations
  - Data storage
  - Log buffers
  - Index
  - Etc...



#### 02 Background & Motivation



# Background

- Redis
  - SET command process
    - 1. Redis client send key-value data to server
    - 2. Redis server check the requested parameters
    - 3. Calculate the hash value of the string for the key to determine the location where data will be stored
    - 4. Make dictEntry & pointing requested key-value data
    - 5. Store dictEntry in dict
  - Additional memory required for data management
    - ✓ dictEntry
      - ✓ Key pointer
      - ✓ Redis object pointer
      - ✓ Next entry pointer
    - ✓ Redis object
      - ✓ Data type
      - ✓ Reference count
      - ✓ Time
  - Redis stores requested data in its original form
    - ✓ Fast data processing performance
    - ✓ Increased burden of data storage & system operation
  - If Redis runs out of available memory...
    - 1. Redis will not be able to execute commands
    - 2. Out-Of-Memory → System Shutdown





## Background

- Redis Database(RDB)
- RDB Logging



- Generate point-in-time snapshots(compressed binary format) of a dataset at specific intervals
- Advantage
  - Small file size
  - Higher performance than AOF
  - Fast recovery speed
- Disadvantage
  - Can not guarantee data persistency





- Redis restarts, it uses RDB file to recover data
- RDB recovery is sequentially performed for all data recorded in the RDB file
- Compared AOF Recovery
  - RDB recovery is faster
  - RDB recovers data by performing only one recovery operation on a key

#### **Motivation**





- Redis utilizes copy-on-write technique to perform RDB
- Copy-on-write: copy the page and reflect changes made to the copied page
- Original page of the copied page is maintained until RDB is finished → Memory occupancy
- Copy-on-write allow Redis to handle requested commands without interruption
- But!!! It causes system memory shortage

#### **Motivation**





- Simulate workloads of data insertion continues to be requested by users
   → The effect of the copy-on-write can be identified
  - Workload A
    - Key 16B, Value 1KB
    - 5,000,000 SET Commands
  - Workload B
    - Key 16B, Value 5KB
    - 5,000,000 SET Commands

#### **Motivation**





- During both workloads, RDB occurs three times each at 60 seconds intervals.
- As the size of the dataset stored in Redis increases, the size of the page copied during the snapshot creation operation increases.
- The longer the snapshot creation task increases, the longer the original pages exist
- The experimental results indicate a significant memory load on Redis used in production environments where a large amount of data is constantly requested.





Data Compression Storage



Compressed value data lookup

- Compress & store value data over 5KB
- To minimize data processing due to compression, use LZF (Lempel-Ziv-Free) compression algorithms
- LZF algorithm: Fast compression speed with low memory consumption
- To maintain data processing performance, key string is stored in the original form without compressing.
- If the value data is stored in compressed form when the GET command is executed, the value data is decompressed and returned.



• Parallel RDB Generation



- Increase the size and number of storable data → Increase logging overhead
- RDB operation is performed using parallel processing
- During RDB operations, log records are written directly to the log file without compressing data
- Data stored in Redis is partitioned into PRDB files and stored





- All PRDB files that exist on disk are used to recover 100% of the data sets stored in Redis
- Performs a recovery by reading the PRDB files that exist on the disk in sequence of file numbers
- Recover data directly to the disk without decompression, even if the value data of the log record recorded in the PRDB file is compressed





• Configuration

Hardware Setup						
СРО	Intel(R) Xeon(R) CPU E5-2660 v2 @ 2.20GHz					
RAM	DDR3 64 GB					
Disk(SSD)	Crucial_CT250MX200 SSD1 250 GB * 3					
Software Setup						
OS	Cent OS 7.3.1611 (Core)					
OS Kernel	3.10.0-514.26.2.el7.x86_64					
Redis version	4.0.10					
Maxmemory option	60 GB					
RDB option	save 60 10000					
Memtier-Benchmark Version	1.2.13					

• Data insertion workload

Workload Name	Key Size (B)	Number of SET Requests	Value size (KB)	Total Data Size	
W1		17388738	0.002	1.72 MB	
W2			0.5	50.35 MB	
W3	12.23	100,000	1	99.18 MB	
W4	W4		10	978.09 MB	
W5			20	1.91 GB	
W6			0.002	8.58 MB	
W7			0.5	251.77 MB	
W8	16	500,000	1	495.91 MB	
W9			10	4.78 GB	
W10			20	9.54 GB	
W11			0.002	42.92 MB	
W12			0.5	1.23 GB	
W13		2,500,000	1	2.42 GB	
W14			10	23.88 GB	
W15			20	47.72 GB	



#### • Configuration

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#### Data lookup workload

Workload Name	Key Size (B)	Number of SET Requests	Value size (KB)	Number of GET Requests
R1				500,000
R2			0.002	1,000,000
R3				1,500,000
R4				500,000
R5			0.5	1,000,000
R6				1,500,000
R7				500,000
R8	16	1,500,000	1	1,000,000
R9				1,500,000
R10			10	500,000
R11				1,000,000
R12				1,500,000
R13				500,000
R14			20	1,000,000
R15				1,500,000



1,000,000 1,500,000

20KB



**Data Insertion Performance** 







**LZF** Compression and Decompression Time





Number	Value	Re	dis	CP-Redis		
of Requests	size (KB)	Copied page (MB)	Time (sec)	Copied page (MB)	Time (sec)	
a Luz	0.002	28	1.10	22	0.44	
	0.5 00 1	30	2.80	36	0.66	
2,500,000		36	3.48	36	0.79	
	10	94	12.09	28	1.59	
Such In		26	26.72	38	2.97	

Copy-on-write overhead due to snapshot creation

					Redi	is 🔳	CP-R	edis				
TIME (SEC)	0.100	0.442 10.439	<b>1.825</b>	0.270 0.124	1.270 0.552	<b></b> 4.998 <b>2.277</b>	0.414	1.685 1.0.570	<b>2.408</b>	■ 2.407 0.165	11.542	57.312
	200,000	1,000,000	5,000,000	200,000	1,000,000	5,000,000	200,000	1,000,000	5,000,000	200,000	1,000,000	5,000,000
	0	.002K	в		0.5KB			1KB			10KB	
					PRDE	3 Rec	overv	Time	e			

